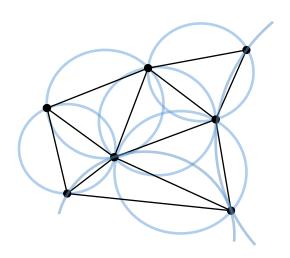
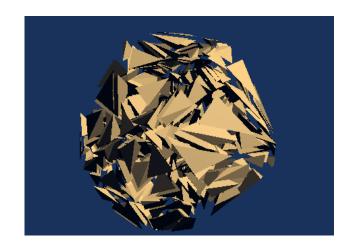
# 3D Delaunay Triangulation

Libin Lu, Weijing Liu, Zhuoheng Yang

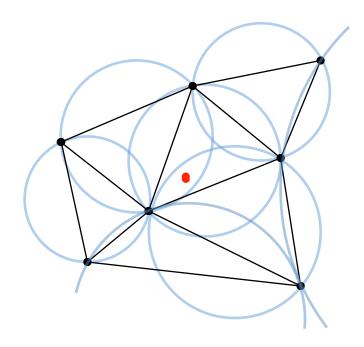
## **Delaunay Triangulation**

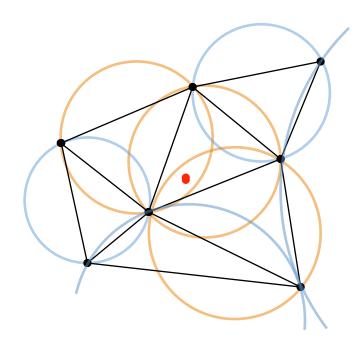


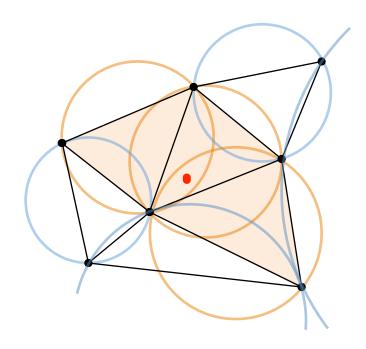


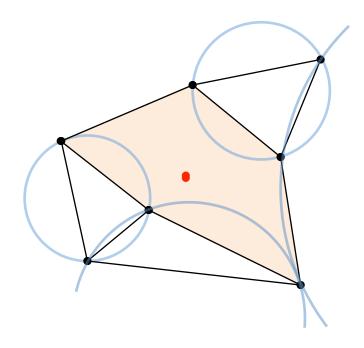
In 2-D, a Delaunay triangulation for a set P of points in a plane is a triangulation DT(P) such that no point in P is inside the circumcircle of any triangle in DT(P).<sup>[1]</sup>

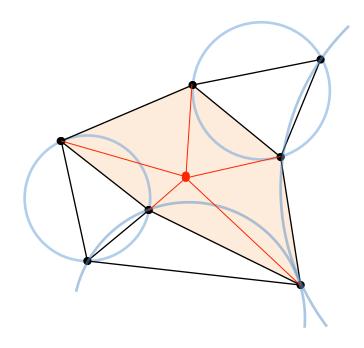
In 3-D, a Delaunay triangulation for a set P of points in space is a triangulation DT(P) such that no point in P is inside the circumsphere of any tetrahedron in DT(P).



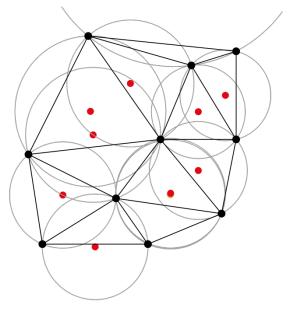




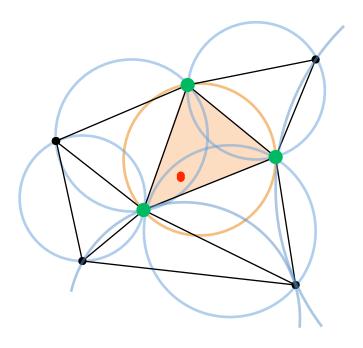




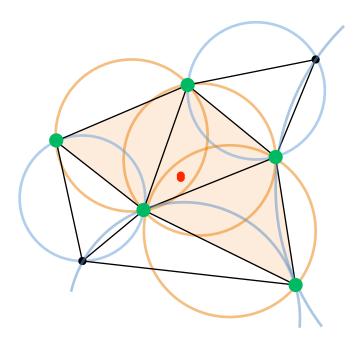
#### **Parallelism**



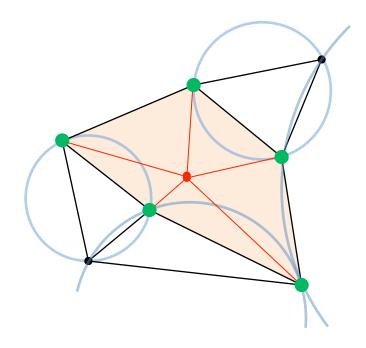
- 1. Each task represents a triangle in plane containing some points.
- 2. Each thread tries to get a task and insert points, triangles far from each other can perform insertion at the same time.



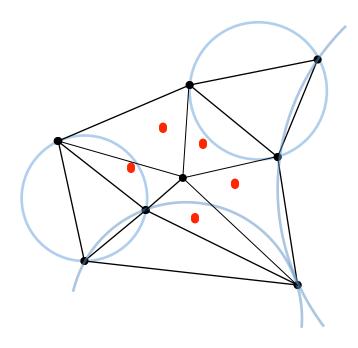
Lock the vertex of task triangle.



Lock more points, all the vertex of the triangles whose circumcircle contain the point to insert.

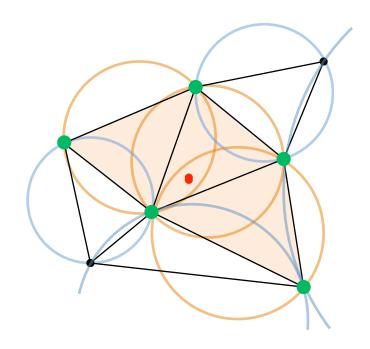


Manipulate them, create new triangles



Sort remaining points into new triangles, put them into task queue, release the locks.

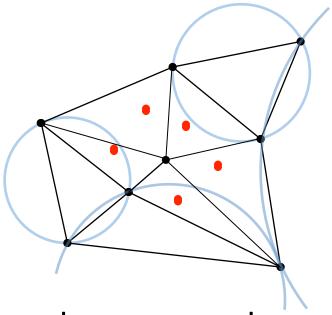
#### **Prevent Deadlocks**



#### Dead lock?

- If lock fail, try again later!
- Why not lock vertex in indexing order?

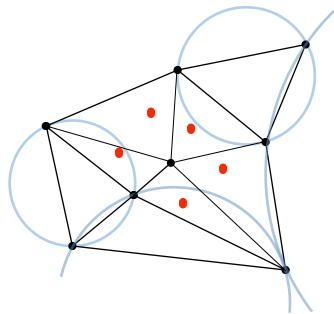
# Using TBB



Tree style relations between tasks

- Not really

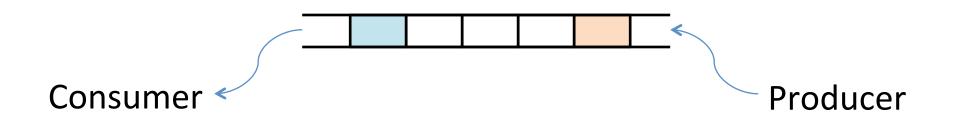
### Using TBB



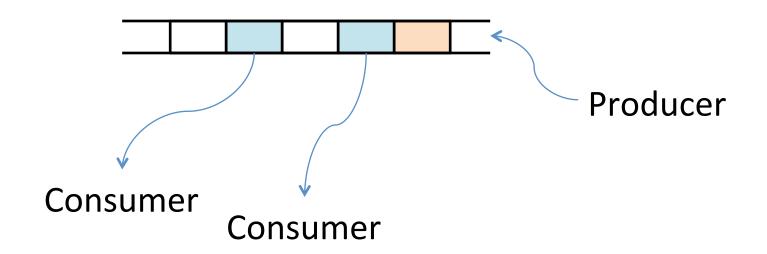
Tree style relations between tasks

- Not really

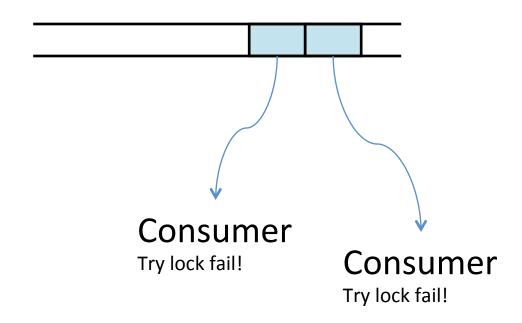
Use customized task queue, draw task randomly from queue.



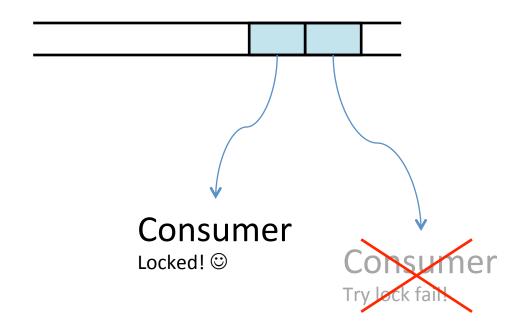
Producer-consumer problem



Fetch a task randomly



Too many conflicts if tasks are highly related. **Starvation!** 



Only allow one consumer if there are too few tasks.

## **Memory Locality**

#### Goal:

Data associated with an insertion task should shares a small number of cache lines, which improves data locality.

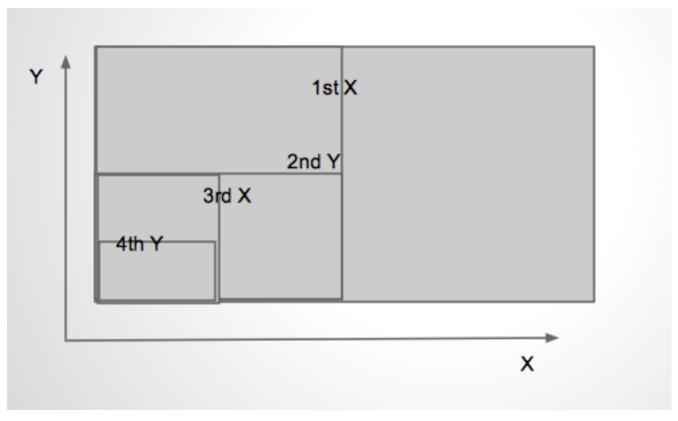
#### Ideas:

- 1. Vertices that are spatially related should also stay close in memory.(done)
- 2. Maybe compress data(not yet)

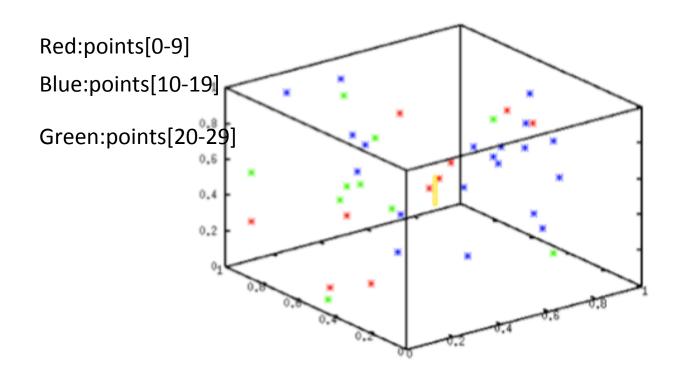
#### **Algorithm: X-Y-Z Cuts**

- 1. Find which of the x,y,z axes has the greatest diameter.
- 2. Find the approximate median(M) of D.
- 3. Partitions the points using M.
- 4. Recursively apply X-Y-Z Cuts to one side of the points first, then the other.
- (Sounds like Quicksort).

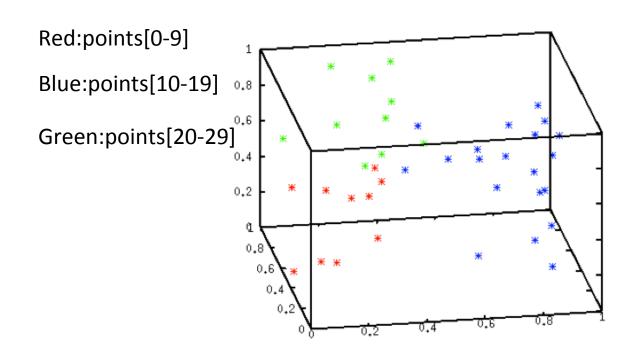
#### Algorithm: X-Y-Z Cuts (Example in 2D)



#### X-Y-Z Cuts (Before sorting)

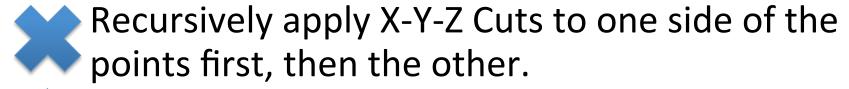


#### X-Y-Z Cuts (After sorting)



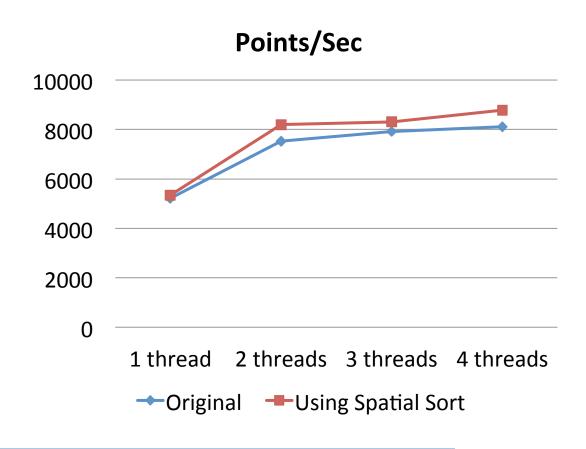
#### X-Y-Z Cuts Optimization

- 1. Inplace sorting. (minimize temporary memory allocation)
- 2. Parallel sorting.
- **Key**: at step #4.



Recursively apply X-Y-Z Cuts to **both** sides of the points **concurrently**.

## Performance With Spatial Sort



#### Conclusion

Difficult to parallel 3D Delaunay triangulation

- Because of complicated data dependencies

More difficult to achieve high efficiency

- hard to distribute irrelative tasks to different threads