

# Integral Equations and Fast Algorithms

## Lecture 25: Singular quadrature, Intro Fast Algorithms

CS 598AK · November 19, 2013

# Outline

Singular quadrature

- Special-purpose methods

- Quadrature by expansion

- QBX method design

Fast Algorithms

# Outline

Singular quadrature

Special-purpose methods

Quadrature by expansion

QBX method design

Fast Algorithms

# Kussmaul/Martensen (aka “Kress”)

Theorem (A special integral [Kress LIE Lemma 8.21])

$$\frac{1}{2\pi} \int_0^{2\pi} \log \left( 4 \sin^2 \frac{t}{2} \right) e^{imt} dt = \begin{cases} 0 & m = 0, \\ -\frac{1}{|m|} & m = \pm 1, \pm 2, \dots \end{cases}$$

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Exciting?

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# KM quadrature demo

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DOF choice for KM...?

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Describe scheme



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Bigger idea hiding in KM...?

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Bigger idea hiding in KM...?

What if you knew how to integrate  
Laplace and wanted to do Helmholtz?

# Singularity subtraction

$$\begin{aligned} & \int \langle \text{Thing } X \text{ you would like to integrate} \rangle \\ &= \int \langle \text{Thing } Y \text{ you } \textit{can} \text{ integrate} \rangle \\ &+ \int \langle \text{Difference } X - Y \text{ which is easy to integrate (numerically)} \rangle \end{aligned}$$

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Drawback?

# High-Order Corrected Trapezoidal Quadrature

- Conditions for new nodes, weights  
( $\rightarrow$  linear algebraic system, dep. on  $n$ )  
to integrate

$$\langle \text{smooth} \rangle \cdot \langle \text{singular} \rangle + \langle \text{smooth} \rangle$$

- Allowed singularities:  $|x|^\lambda$  (for  $|\lambda| < 1$ ),  $\log|x|$
- Generic nodes and weights for log singularity
- Nodes and weights copy-and-pasteable from paper

[Kapur, Rokhlin '97]

# High-Order Corrected Trapezoidal Quadrature

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$\langle sm$

- Allowed singularities
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- Nodes and weights

[Kapur, Rokhlin '97]

[Alpert '99] conceptually similar:

- Hybrid Gauss-Trapezoidal
- Positive weights
- Somewhat more accurate (empirically) than K-R
- Similar allowed singularities ( $\lambda > -1$ )
- Copy-paste weights

# Generalized Gaussian

- “Gaussian” :
  - Integrates  $2n$  functions exactly with  $n$  nodes
  - Positive weights
- Clarify assumptions on system of functions (“Chebyshev system”) for which Gaussian quadratures exist
- When do (left/right) singular vectors of integral operators give rise to Chebyshev systems?
  - In many practical cases!
- Find nodes/weights by Newton’s method
  - With special starting point
- Very accurate
- Nodes and weights for download

[Yarvin/Rokhlin ‘98]

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Generalizes to  $nD \dots$



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- Nodes and weights

Generalizes to  $nD$ ...

...if you know how to make Newton’s method converge

[Yarvin/Rokhlin ‘98]

# Singularity cancellation: Polar coordinate transform

$$\begin{aligned} & \iint_{\partial\Omega} K(\mathbf{x}, \mathbf{y}) \varphi(\mathbf{y}) ds_{\mathbf{y}} \\ &= \\ & \int_0^R \int_{\mathbf{x}+\mathbf{r} \in \partial\Omega \cap \partial B(\mathbf{x}, r)} K(\mathbf{x}, \mathbf{x} + \mathbf{r}) \varphi(\mathbf{x} + \mathbf{r}) d\langle \text{angles} \rangle r dr \\ &= \\ & \int_0^R \int_{\mathbf{x}+\mathbf{r} \in \partial\Omega \cap \partial B(\mathbf{x}, r)} \frac{K_{\text{less singular}}(\mathbf{x}, \mathbf{x} + \mathbf{r})}{r} \varphi(\mathbf{x} + \mathbf{r}) d\langle \text{angles} \rangle r dr \end{aligned}$$

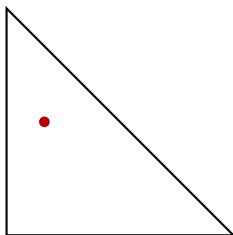
where  $K_{\text{less singular}} = K \cdot r$

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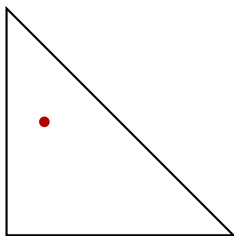
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# Quadrature on triangles



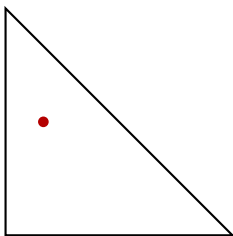
**Problem:** Singularity can sit *anywhere* in triangle

# Quadrature on triangles



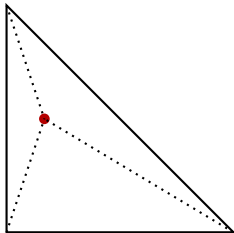
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→ need *lots* of quadrature rules (one per target)

# Quadrature on triangles



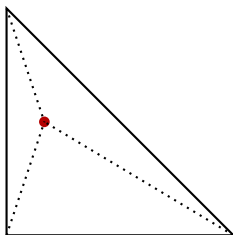
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→ need *lots* of quadrature rules (one per target) ...!?

Possible issue?



# Kernel Regularization

Singularity makes integration troublesome: *Get rid of it!*

$$\frac{\dots}{\sqrt{(x-y)^2}} \rightarrow \frac{\dots}{\sqrt{(x-y)^2 + \varepsilon^2}}$$

Use Richardson extrapolation to recover limit as  $\varepsilon \rightarrow 0$ .

(May also use geometric motivation: limit along line towards singular point.)

Primary drawbacks:

- Low-order accurate
- Need to make  $\varepsilon$  smaller (i.e. kernel more singular) to get better accuracy

# Kernel Regularization

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(May also use geometric motivation: limit along line towards singular point.)

Primary drawbacks:

- Low-order accurate
- Need to make  $\varepsilon$  small for better accuracy

Can take many forms—for example:

- Convolve integrand to smooth it ( $\rightarrow$  remove/weaken singularity)
- Extrapolate towards no smoothing

Related: [Beale/Lai '01]

# Outline

Singular quadrature

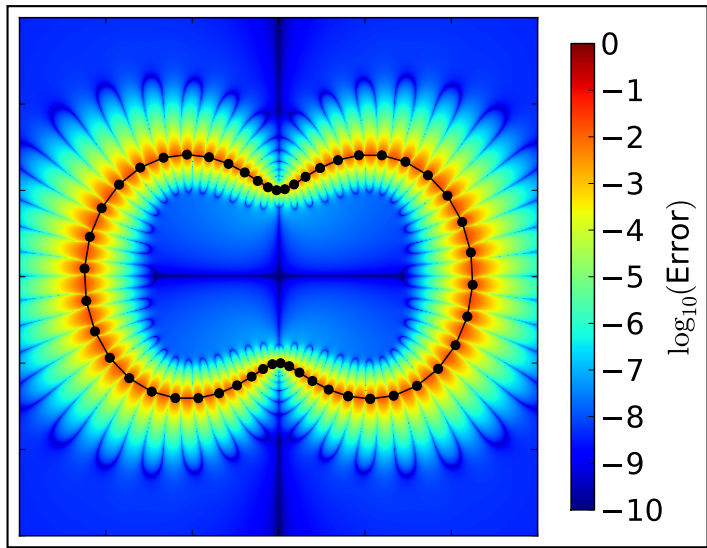
Special-purpose methods

**Quadrature by expansion**

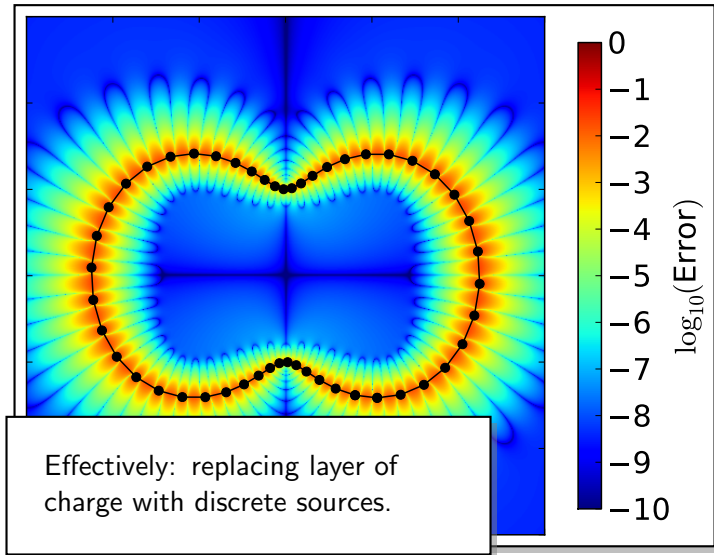
QBX method design

Fast Algorithms

## Using just the trapezoidal rule

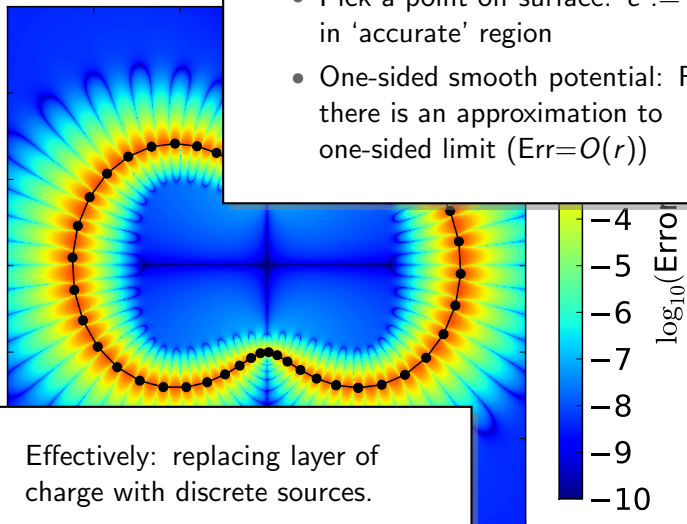


## Using just the trapezoidal rule



**Idea:**

- Pick a point off surface:  $c := x + \hat{n}r$  in 'accurate' region
- One-sided smooth potential: Field value there is an approximation to one-sided limit ( $\text{Err} = O(r)$ )

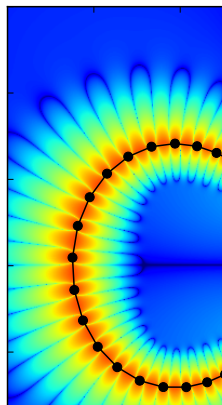


Effectively: replacing layer of charge with discrete sources.

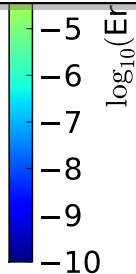
**Idea:**

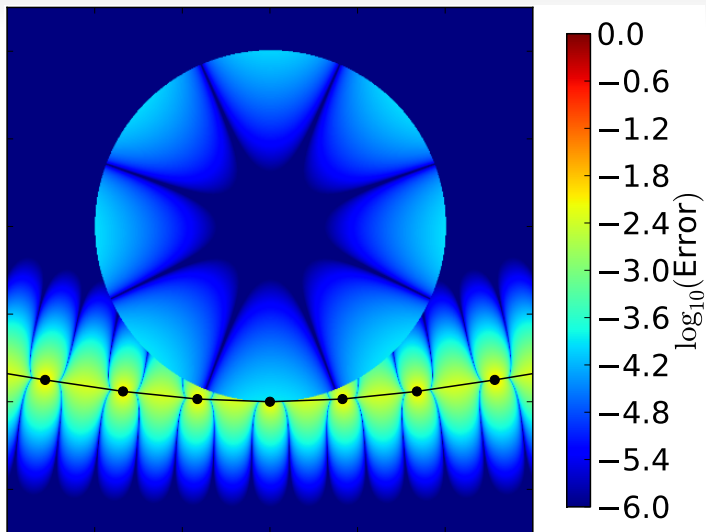
- Pick a point off surface:  $c := x + \hat{n}r$  in 'accurate' region
- One-sided smooth potential: Field value there is an approximation to one-sided limit ( $\text{Err} = O(r)$ )

*But:* Can do much better!

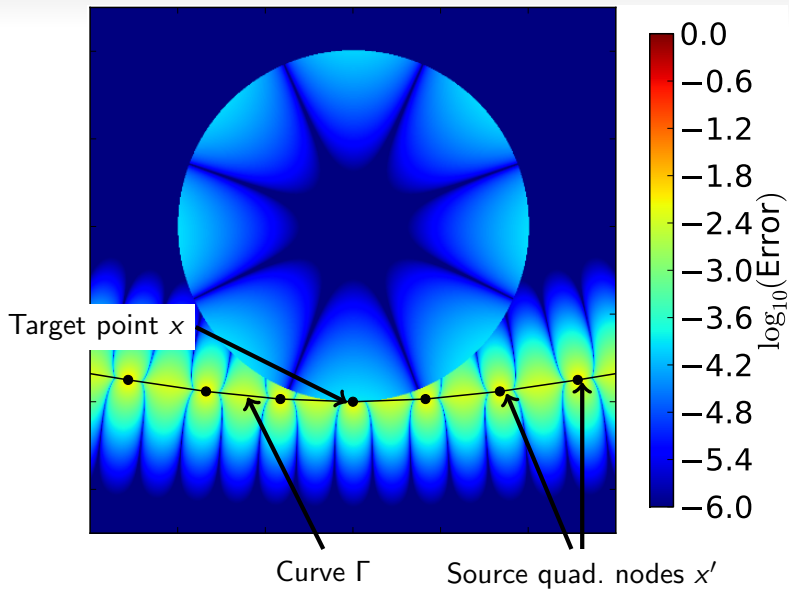


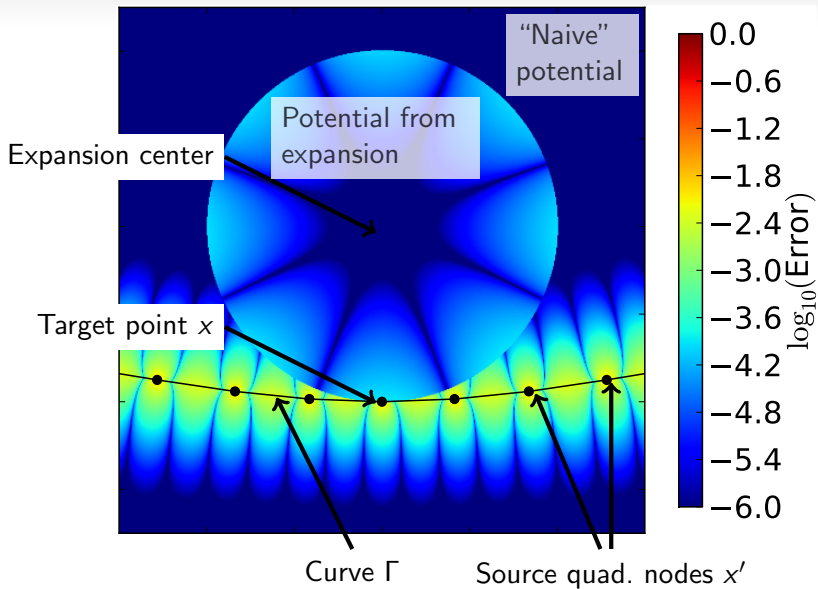
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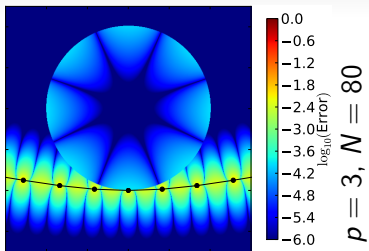


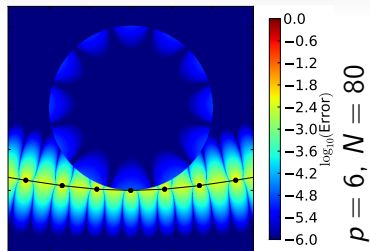
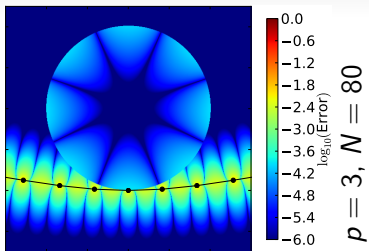


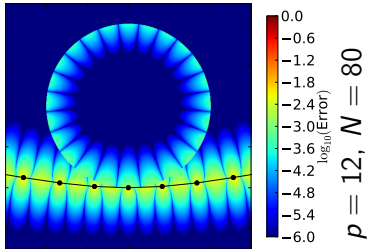
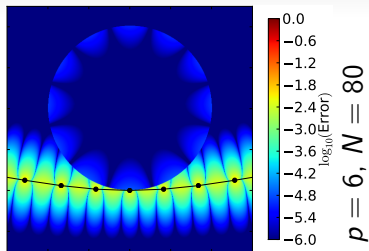
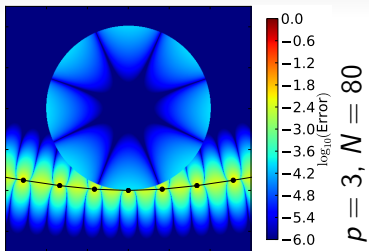


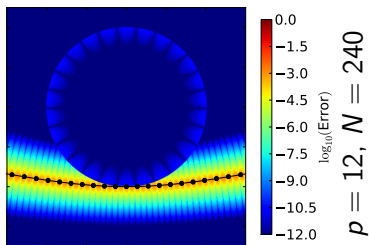
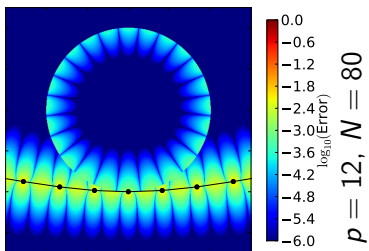
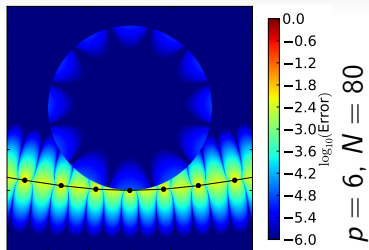
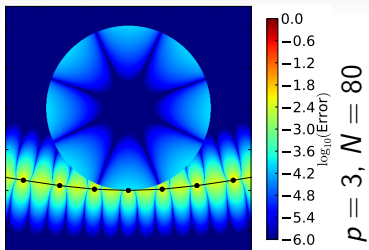






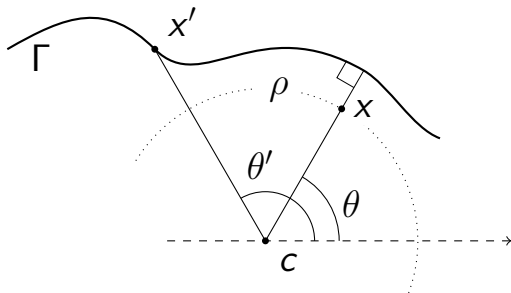






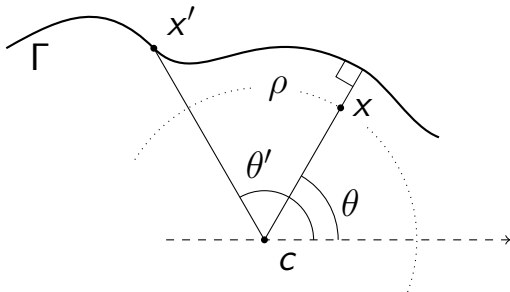
# QBX in formulas: Notation, Basics

## Graf's addition theorem



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## Graf's addition theorem



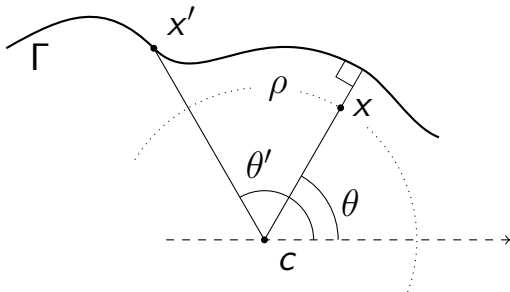
$$H_0^{(1)}(k|x - x'|) = \sum_{l=-\infty}^{\infty} H_l^{(1)}(k|x' - c|) e^{il\theta'} J_l(k|x - c|) e^{-il\theta}$$



QB

Requires:  $|x - c| < |x' - c|$  ("local expansion")

Graf's addition theorem



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# QBX in formulas: Formulation, discretization

Compute layer potential on the disk as

$$S_k \sigma(x) = \sum_{l=-\infty}^{\infty} \alpha_l J_l(k\rho) e^{-il\theta}$$

with

$$\alpha_l = \frac{i}{4} \int_{\Gamma} H_l^{(1)}(k|x' - c|) e^{il\theta'} \sigma(x') dx' \quad (l = -\infty, \dots, \infty)$$

$S\sigma$  is a smooth function *up to*  $\Gamma$ .

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Two limits ( $p, N \rightarrow \infty$ )! Experiment showed: *order matters!*

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Two limits ( $p, N \rightarrow \infty$ )! Experiment showed: *order matters!*

*And:* failure and repair not actually surprising.

# Error result

$$\left| S\sigma(x) - \sum_{l=-p}^p \alpha_l^{\text{QBX}} J_l(k|x-c|) e^{-il\theta_{cx}} \right|$$
$$\leq \left( \underbrace{C_{p,\beta} r^{p+1} \|\sigma\|_{\mathcal{C}^{p,\beta}(\Gamma)}}_{\text{Truncation error}} + \underbrace{\tilde{C}_{p,2q,\beta} \left(\frac{h}{4r}\right)^{2q} \|\sigma\|_{\mathcal{C}^{2q,\beta}(\Gamma)}}_{\text{Quadrature error}} \right)$$

Proof sketch:

1. First, assume exact calculation of coefficients
2. Estimate tail of expansion
3. Estimate quadrature error in coefficients (derivatives/...)
4. Sum quadrature errors in truncated expansion

[K, Barnett, Greengard, O'Neil '12 (submitted)]



# Outline

## Singular quadrature

- Special-purpose methods

- Quadrature by expansion

- QBX method design**

## Fast Algorithms

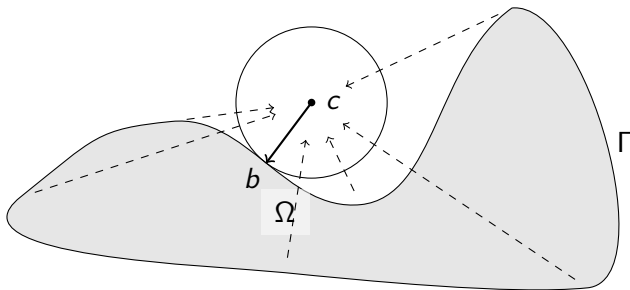
# Achieving high order

$$\text{Error} \leq \left( C \underbrace{r^{p+1}}_{\text{Truncation error}} + C \underbrace{\left(\frac{h}{r}\right)^q}_{\text{Quadrature error}} \right) \|\sigma\|$$

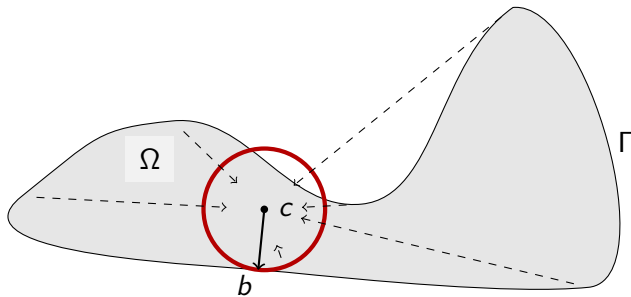
Two approaches:

- *Asymptotically convergent:*  $r = \sqrt{h}$ 
  - Error  $\rightarrow 0$  as  $h \rightarrow 0$
  - ➖ Low order:  $h^{(p+1)/2}$
- *Convergent with controlled precision:*  $r = 5h$ 
  - ➖ Error  $\not\rightarrow 0$  as  $h \rightarrow 0$
  - High order:  $h^{p+1}$   
to controlled precision  $\varepsilon := (1/5)^q$

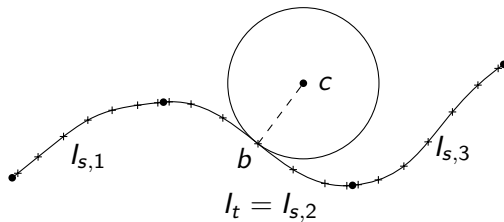
# “Global” QBX: Dealing with geometry



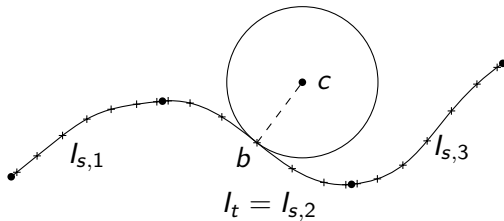
## “Global” QBX, part II



# “Local” QBX

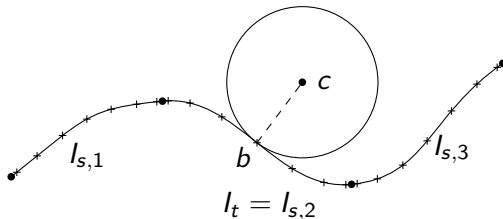


# “Local” QBX



Makes geometry processing much simpler

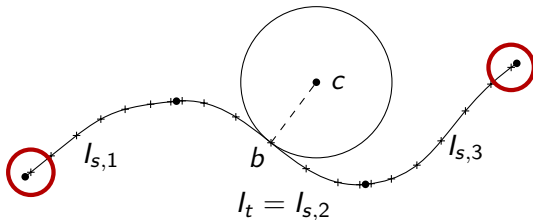
# "Local" QBX



**Problem:** Expanded field becomes non-smooth (because of end singularities)

...kes geometry process-  
much simpler

# "Local" QBX

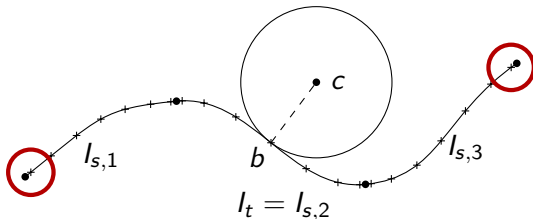


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# “Local” QBX



**Problem:** Expanded field becomes non-smooth (because of end singularities)

**Idea:** Manage as additional, finite error contribution (using  $p$ ,  $h \propto r$ )

...kes geometry process-  
much simpler

## Other layer potentials

Can't just do single-layer potentials:

$$\alpha_l^D = \frac{i}{4} \int_{\Gamma} \frac{\partial}{\partial \hat{n}_{x'}} H_l^{(1)}(k|x' - c|) e^{il\theta'} \mu(x') dx'.$$

Even easier for target derivatives ( $S'$  et al.):

Take derivative of local expansion.

**Analysis says:** Will lose an order.

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**Slight issue:** QBX computes one-sided limits.

Fortunately: Jump relations are known—e.g.

$$(PV)D^* \mu(x)|_{\Gamma} = \lim_{x^{\pm} \rightarrow x} D\mu(x^{\pm}) \mp \frac{1}{2} \mu(x).$$

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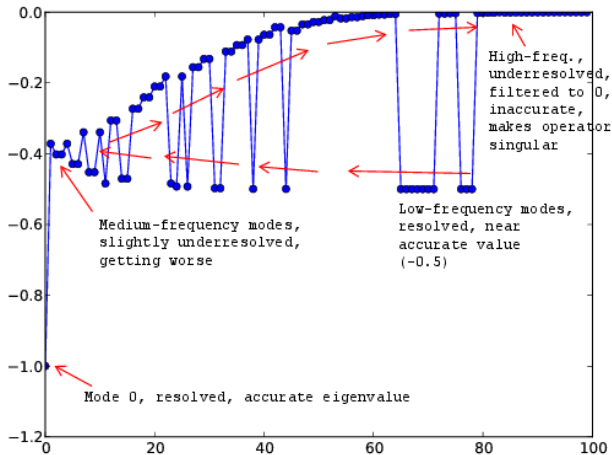
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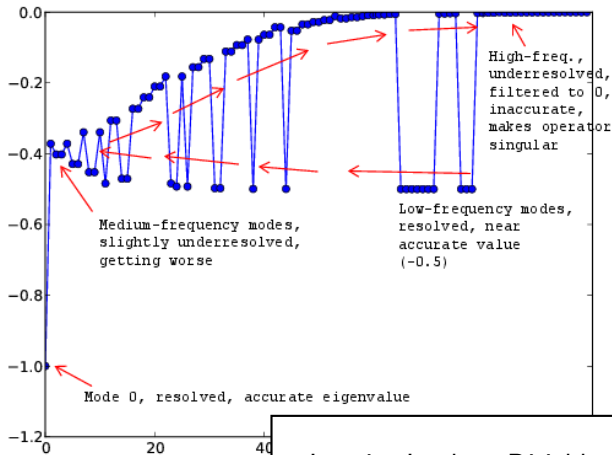
$$(PV)D^*\mu(x)|_{\Gamma} = \lim_{\epsilon \rightarrow 0} \frac{D\mu(x^{\pm})}{2} \mp \frac{1}{2}\mu(x)$$

*Alternative:* Two-sided average

# Spectral behavior



# Spectral behavior



Interior Laplace Dirichlet problem would try to invert this operator.

## Spectral behavior, part II

- QBX *wants* to approximate a compact operator—let it:

$$D\mu(x) = \frac{1}{2} \left( \lim_{x^+ \rightarrow x} D\mu(x^+) + \lim_{x^- \rightarrow x} D\mu(x^-) \right).$$

Simply use two QBX applications.

- *Predictably benign spectral behavior* at high frequencies.

Important for iterative solvers (e.g. GMRES)

Not many competing schemes have that!

# Outline

Singular quadrature

Fast Algorithms



# Integral equations + computers

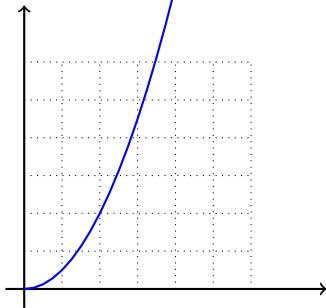
$$O(n^2)$$

# Integral equations: computational expense

Why is  $O(n^2)$  a problem?

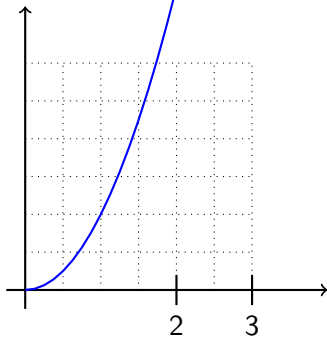
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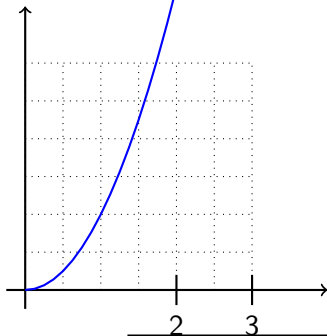
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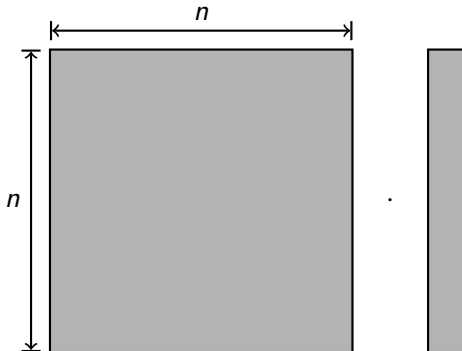


Why is  $O(n^2)$  *storage complexity* a problem specifically?

# Integral equations: computational expense

$O(n^2)$  not in principle incorrect:

Natural complexity of a dense mat-vec.



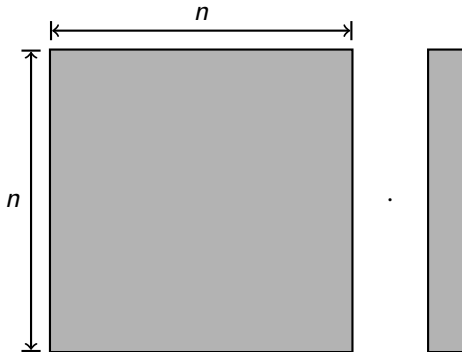
**Or**, to be more precise:

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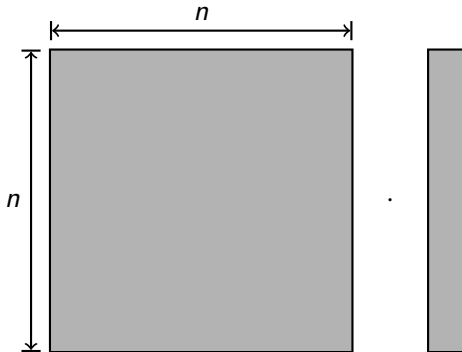
**Or**, to be more precise:

$$O(\langle \# \text{sources} \rangle \cdot \langle \# \text{targets} \rangle) = O(\langle \# \text{rows} \rangle \cdot \langle \# \text{columns} \rangle)$$

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Special dense matrices with  
faster mat-vecs?



# Faster dense mat-vecs

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Define numerical rank

# Questions?

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